## Lecture 14 - ARQ and Flow Control

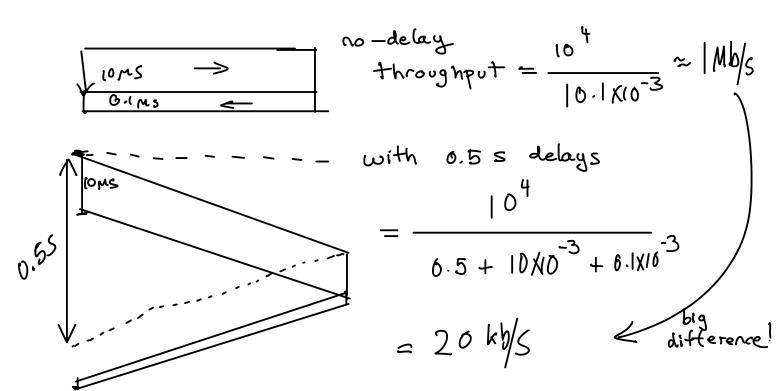
**Exercise 1**: Create a table summarizing the different types of ARQ. Include: throughput, transmitter memory, receiver memory and relative complexity.

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stop& bait		1	1	L
go-back N	Н	N	1	M
selective repeat	H	N	N	H

**Exercise 2**: A data communication system operates at 1 Mb/s and uses 10,000-bit data frames and 100-bit ACK frames. What are the frame durations? What is the throughput if there is no channel delay and no errors? If the round-trip channel delay is a 0.5s (typical for satellite links)? If go-back-N ARQ is used, assuming the transmitter can store 200 unacknowledged frames?

$$\frac{10^{4}}{10^{6}} = 10^{-2} = 10^{-2}$$

$$\frac{10^{2}}{10^{6}} = 1 \times 10^{-4} = 0.1 \text{ms}$$



w/ go-book N (kno errors):

if can transmit for 200 x loms = 2s

before stopping for an ACK, we will

receive the ACK before need to stop

transmitting & throughput will be = 1Mb/s.

**Exercise 3**: A communication system loses every 10th frame (e.g. due to periodic noise bursts). Ignoring ACK overhead, what is the throughput using go-back-N ARQ? Using Selective ARQ?

Answer depends on delay

If no delay, some throughput (  $\frac{10 \text{ received}}{11 + 120 \text{ smissions}} \approx 91\%$ )

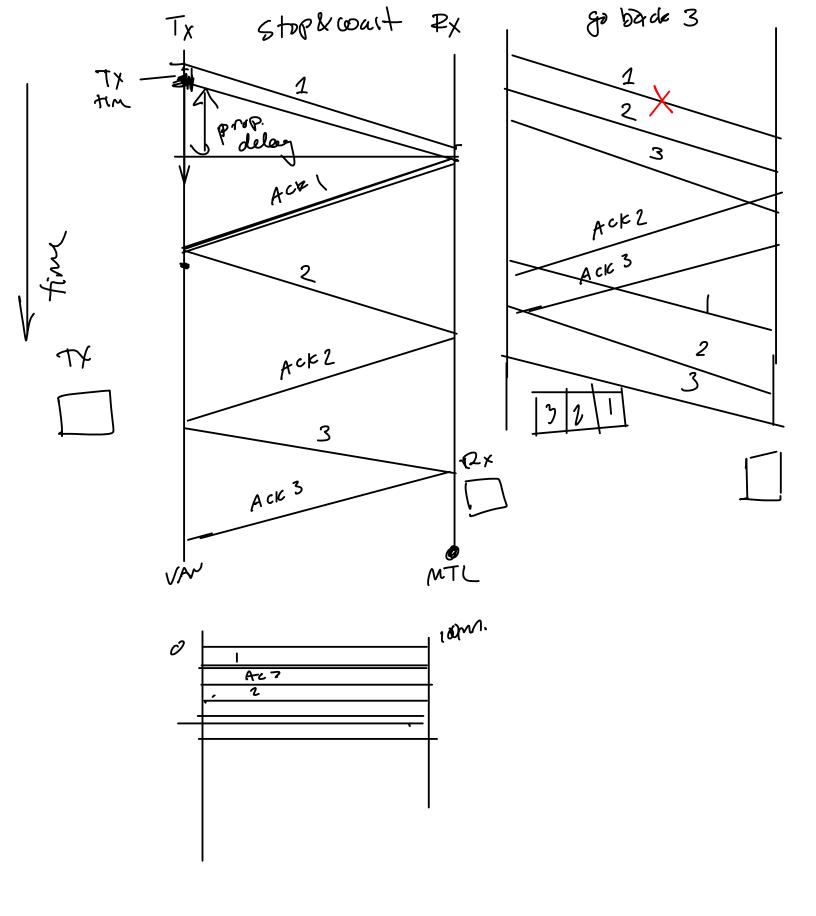
If have to repeat N frames due

to delay then throughput is  $\frac{10 \text{ received}}{10 + 1 + N}$ e.g. N=5 throughput =  $\frac{10}{16} \approx 63\%$  for go-back-N

but still  $\frac{10}{11} = 91\%$  for selective repeat.

**Exercise 4**: Which of the above flow control methods can be used with frame-oriented protocols? On unidirectional links?

- frame oriented protocds can delay ACKs to do flow control (can also use other methods).
- unidirectional links cannot use software flow control (or ACK packets), only how f/c.



Tx buffer selective repeat

| 3 2 1 |

Ack 2

Ack 3

Ack 1

2x 50+4=1