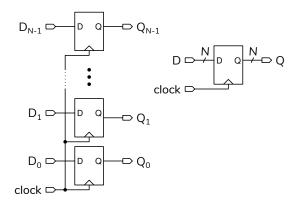
State Machines

This lecture defines state machines and describes how to document them and how to implement them using Verilog. After this lecture you should be able to: design a state machine based on an informal description of its operation, document it using state transition diagrams and tables, write a synthesizable Verilog description of it and convert between these three descriptions.

Introduction

Registers

We can connect *N* flip-flops to the same clock to form an *N*-bit *register*. The common clock loads every flip-flop at the same time:

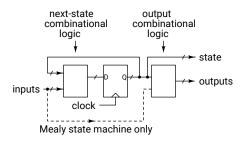


The Verilog always_ff statement creates a register.

State Machines

The *state* of a register is its value. A *state machine* is a description of how the state changes.

A state machine is implemented as a register whose value on the next rising edge of the clock is computed from the current state and, optionally, inputs:



The state register is loaded on each rising edge of the clock, but if it's loaded with its current value then there is no change of state.

The output of a state machine is a function of its state. In some cases the output may be the state itself.

The above describes a *Moore* state machine. A *Mealy* state machine is one where the output is a function of the current state as well as the inputs.

State Machine Descriptions

For example, consider a state machine with two bits of state that sequences through the values 00, 01, 10, 11 and back to 00. The output should be 0 in states 00 and 01 and 1 in states 10 and 11. There is an input named reset. The state is set to 00 if reset is 1.

Truth Tables

We can describe a state machine using truth-tables. One table has columns for the current state, the input value(s), and the corresponding next state. Another table has columns for the current state and the output.

We could write the state transition table and the output table as:

state	reset	next
		state
00	0	01
01	0	10
10	0	11
11	0	00
00	1	00
01	1	00
10	1	00
11	1	00

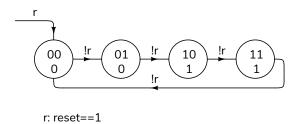
state	output
00	0
01	0
10	1
11	1

Exercise 1: For the state machine described above, if the current state is 01, what will be the next state? When will the state change? What is the output in state 00? In state 01? In state 10?

State Transition Diagrams

A state machine can also be described by a state transition diagram drawn using the following conventions: (1) each state is represented by a circle labelled with the state name or value; (2) each state shows the output for that state (unless the state is also the output value); (3) arrows show possible transitions between states; (4) transitions happen when the expression labelling an arrow is non-zero (true)¹; (5) unlabelled transitions happen unconditionally; (6) transitions with no origin come from every state; (7) conditions that don't cause a state change are not shown²; (8) only one transition out of a state may be true (the expressions on the arrows leaving each state must be mutually exclusive).

The following is a state transition diagram for the state machine above:



Exercise 2: Modify the diagram so the state machine counts to 11 and stops.

Exercise 3: Add a down input that cause the values to count down.

Simplifications

We can simplify truth tables using the following conventions: (1) x can be used for "don't care" in state or input columns; (2) the first matching row applies; (3) if there is no matching row there is no change of state; (4) expressions can be used to define the next state as a function of the current state and the input.

For example, we could rewrite the above state transition table as:

state	reset	next
		state
XX	1	00
11	0	00
n	0	n+1

 $^{^{1}}$ Abbreviations for these expressions can be used to keep the diagram tidy.

Exercise 4: What will be the next state if the state is 00 and the reset input is 1? If the state is 00 and the reset input is 0? When does the state change? When does reset affect the output?

States can be labelled with names instead of numerical values.

Exercise 5: Write the above state transition and output tables using state names A, B, C, and D.

State Machines in Verilog

// 2-bit clock divider with reset

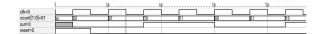
A state machine can be written in Verilog using one conditional operator for each table row (or for each diagram arrow). The condition is expression that labels the arrow in the diagram or an expression that is true for the current state and inputs for that transition. The true value is the next state for that transition. The false value is the next conditional operator or, if no transition matches, the current state.

A straightforward translation of the truth tables above would be written as:

```
module ex79
   ( output logic [1:0] count,
     output logic out,
     input logic reset, clk );
   always_ff @(posedge clk) count
      <= reset && count == 2'b00 ? 2'b00 :</pre>
          reset && count == 2'b01 ? 2'b00 :
          reset && count == 2'b10 ? 2'b00 :
          reset && count == 2'b11 ? 2'b00 :
         !reset && count == 2'b00 ? 2'b01 :
         !reset && count == 2'b01 ? 2'b10 :
         !reset && count == 2'b10 ? 2'b11 :
         !reset && count == 2'b11 ? 2'b00 :
         count ;
   assign out
     = out == 2'b00 ? 1'b0 :
       out == 2'b01 ? 1'b0 :
       out == 2'b10 ? 1'b1 : 1'b1 ;
endmodule
Or, more simply:
// 2-bit clock divider with reset
module ex67
   ( output logic [1:0] count,
     output logic out,
     input logic reset, clk );
   always_ff @(posedge clk) count
      <= reset ? 2'b00 :
         count == 2'b11 ? 2'b00 :
         count + 1'b1;
   assign out = count == 2'b10 || count == 2'b11 ;
```

endmodule

²Some authors do not allow this.



Counters

A counter is a common example of a state machine. Typical inputs include those to restart the sequence (typically called reset), to pause or continue the sequence (hold or enable), or change the order of the values generated (e.g. up/down).

Exercise 6: Show the state transition diagram and table for a 2-bit counter with reset, enable, and down inputs. reset should have priority over enable which should have priority over down. Write the Verilog.

Timers and Clock Dividers

It takes N clock periods for a counter to count down³ from N-1 to 0. If the clock period is T seconds then the time taken is NT seconds. A circuit can thus create a delay of NT seconds by counting N clock cycles⁴ **Exercise 7:** What value of N would result in a 20 ms delay if the clock frequency is 50 MHz? How many bits are needed for this timer's register?

If the counter is reset to N-1 when it reaches 0 then the count values will be periodic with a period NT. If some event happens each time the count reaches a specific value (e.g. 0) then this event happens with period NT (and thus a frequency $\frac{1}{NT}$).

Exercise 8: Assume the timer above is reset to N-1 each time it reaches 0. For how long is the register value 0? What are the period and frequency of a signal that is inverted each time the count reaches 0?

State Encodings

States may be represented ("encoded") in different ways:

Binary States are encoded as binary numbers. This requires the fewest number of flip-flops. It is used for state machines with many states such as counters. In the above example the binary encoding of the states would be **00**, **01**, **10** and **11**.

One-Hot There is one flip-flop for each state. Only one flip-flop at a time can be set to 1. This is used when there are few states because simplifies the next-state and output logic. In the above example a one-hot encoding might be 1000, 0100, 0010 and 0001.

Output Each state is encoded as the output for that state. This eliminates the need for any logic to determine the output as a function of the state. However, this is only possible when the output is different for each state. The above example could not use an output state encoding because the output is 0 for states 00 and 01 and the output is 1 for states 10 and 11.

Input The state encoding is a sequence of previous inputs. This is used when the output can be determined from a small number of previous inputs. For example, a combination lock might unlock after the correct sequence of inputs.

Exercise 9: How many bits need to be considered to detect a specific state when a binary encoding is used? How many need to be considered if a one-hot encoding is used?

Exercise 10: If we used 8-bits of state information, how many states could be represented? What if we used 8 bits of state but used a "one-hot" encoding?

 $^{^3}$ Timers traditionally count down from N-1 to 0 rather than up from 0 to N-1 because it's simple to determine when the count reaches 0: adding -1 does not cause a carry.

 $^{^4}$ The time includes clock cycles during which the counter has values N-1 through 0 inclusive.