## Solutions to Midterm Exam

## Question 1

The two versions of this question used the same expressions arranged in different orders. The following code computes the size and value of each expression:

```
// solutions for midterm exam question 1
module midterm ;
```

    logic [7:0] a = 8'h08 ;
    logic \([3: 0] \mathrm{b}[3: 0]='\{4 ' b x x x x, 4 ' h 1,4 ' d 2,4 ' b 0011\} ;(1,2,4,8)\) as shown below.
    logic signed [7:0] c = 8'b1 ;
    -define ans(expr) \}
\$display("\%16s: bits: \%2d value: \%8b", \}
-"expr`",\$bits (expr), expr) ;
initial begin
-ans ( 1 \&\& 2 ) ;
-ans ( \{4\{c[0]\}\}) ;
ans ( b[0][1:0] ) ;
-ans ( 4'sb1000 >>> 3 ) ;
-ans ( ${ }^{-c}$ ) ;
-ans ( $\{\mathrm{a}[0], \mathrm{b}[0]\}$ ) ;
ans ( c[0] ? b[2] : b[1] ) ;
ans ( $2 * \mathrm{c}$ ) ;
-ans( a|c) ;
end
endmodule
and the output is:

## Question 2

(a) There were two versions of the question. In both the clock signal is initialized to 0 and the simulation runs for six delays of $1 \mu \mathrm{~s}$ ( $6 \mu \mathrm{~s}$ total). There are two rising clock edges (at 1,3 and $5 \mu \mathrm{~s}$ ).
(b) In both versions of the question x is always set to $\mathrm{y}+\mathrm{y}$ in an always_comb block so the values are

```
            1 && 2: bits: }1\mathrm{ value: 00000001
```

            1 && 2: bits: }1\mathrm{ value: 00000001
            {4{c[0]}}: bits: 4 value: 00001111
            {4{c[0]}}: bits: 4 value: 00001111
            b[0][1:0]: bits: 2 value: 00000011
            b[0][1:0]: bits: 2 value: 00000011
    4'sb1000 >>> 3: bits: 4 value: 00001111
    4'sb1000 >>> 3: bits: 4 value: 00001111
            ^c: bits: 1 value: 00000001
            ^c: bits: 1 value: 00000001
            {a[0],b[0]}: bits: 5 value: 00000011
            {a[0],b[0]}: bits: 5 value: 00000011
    c[0] ? b[2]:b[1]: bits: 4 value: 00000001
c[0] ? b[2]:b[1]: bits: 4 value: 00000001
2*c: bits: 32 value: 00000010
2*c: bits: 32 value: 00000010
alc: bits: 8 value: 00001001

```
            alc: bits: 8 value: 00001001
``` the clock signal is initialized to 0 and the simula
tion runs for six delays of \(1 \mu \mathrm{~s}(6 \mu\) stal). There
always twice those of \(y\). In one case the initial value of \(y\) is 2 and so the initial value of \(x\) is 4 . In the one case the initial value of \(y\) is 1 and so the initial value of x is 2 .
(c) The always_ff block executes and updates \(y\) on each rising clock edge. The value of y is set to twice the current value. So in the first version \(y\) is set to \((2,4,8, \mathrm{~h} 10)\); in the second it is set to
; \((1,2,4,8)\) as shown below.

The simulation results are as follows:
```

